

Ad Hoc Network Secure Protocol Simulator



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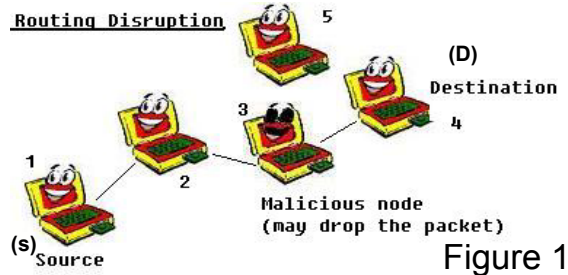
Motivation:

- 1) Simulator can compare the security performance of different protocols
- 2) Ad hoc network can be established without help from a fixed infrastructure
- 3) Security is a critical issue since military work may be involved

Types of attacks:

- 1) Resource consumption attacks (waste bandwidth)
- 2) Routing disruption attacks (drop routing packets)

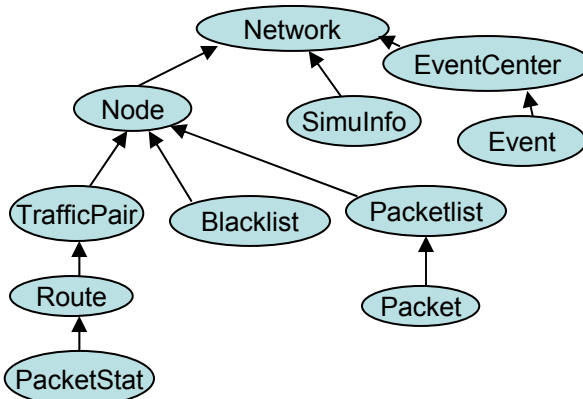
Routing Disruption



Simulator using C++:

- Object-oriented
- Easy to add more components or easy to change the scenario

Program Structure:



- Network**
 - Controls all other objects
 - Contains nodes and graphs to simulate an ad hoc network
 - Processes the events get from EventCenter
- Event**
 - Types of events: generate packet, forward packet, update graph
- EventCenter**
 - Registers events to an array, sort the event order by priority
 - Network class processes events from here
- SimulInfo**
 - Contains all simulation information: good node #, bad node # area of network, transmission range of each node
- Node**
 - Good node** – never lies **Bad node** – may lie to frame others
 - Contains position of itself, moving speed, moving direction, pause time between each move, power level, status of itself, packetlist
 - Most important:** contains Blacklist, TrafficPair
- Packetlist**
 - linked list of packets
- Packet**
 - Can be data packet, route request packet, route reply packet
 - Knows sender & receiver
 - Since DSR is the routing protocol, route is sent along with the packet
- TrafficPair**
 - Knows it's sender and receiver
 - A list of valid routes for this traffic pair

- PacketStat**
 - RN(A,S,Ri), number of packets received by hop A for S via Ri route
 - FN(A,S,Ri), number of packets forward by hop A for S via Ri route
 - P(A,S,Ri) = $\frac{FN(A,S,Ri)}{RN(A,S,Ri)}$ packet delivery ratio of hop A for S via Ri route
 - Example:** In figure 1, node 2 claims it has forwarded all packets to node 3(malicious). But node 3 has not forwarded all packets to node 4
Result: $P(2,S,Ri) = 1$ & $0 < P(3,S,Ri) < 1$
 - Blacklist**
 - It contains the honesty score of other nodes
 - Honesty score: start with a 1 for every node. Score will go down if a node is suspected cheating
 - Example:**
In figure 1, node 2 claims it has forwarded all packets to node 3(malicious). Node 3 tries to frame node 2, claims it receives none
- $FN(2,S,Ri) \neq RN(3,S,Ri)$

Someone lies, but don't know which one lies, Both 2 and 3 are suspected

Honesty Score:
 $H(2,S) = \alpha H(2,S)$
 $H(3,S) = \alpha H(3,S)$

As time goes on I have enough evidence to say 3 cheated. Honesty score of 2 will go back up

Honesty Score:
 $H(2,S) = \frac{H(2,S)}{\alpha^m}$
 $H(3,S) = 0$
 $m = \#$ of framed by 3

Evaluation of routes:

$$Q(R_i) = \prod_{A \in R_i} P(A,S) * H(A,S) - \lambda * L_i$$

According to the information from PacketStat and Blacklist, we can determine which route is more secure

L_i = number of hops in a route

λ = determine how important is the number of hops

**if the traffic is already heavy, shorter route preferred, we weigh L_i more

Future work:

- 1) Finish the implementation
- 2) Additional routing protocol classes