ENEE244-010x Digital Logic Design

Lecture 6

Announcements

- Homework 2 due today
- Homework 3 up on webpage (tonight), due Wed, 9/30.
- Class is canceled on Wed, 9/23
- Substitute on Mon, 9/28

Agenda

- Last time:
 - Manipulations of Boolean Formulas (3.6)
- This time:
 - Gates and Combinational Networks (3.7)
 - Incomplete Boolean Functions and Don't Care Conditions (3.8)
 - Universal Gates (3.9.3)

Digital Logic Gates

- AND f(x,y) = xy
- OR f(x, y) = x + y
- NOT (Inverter) $f(x) = \overline{x}$
- Buffer (Transfer) f(x) = x
- NAND $f(x,y) = \overline{(xy)} = \overline{x} + \overline{y}$
- NOR $f(x,y) = \overline{x+y} = \overline{x} \, \overline{y}$
- XOR $f(x, y) = x \overline{y} + \overline{x} y = x \oplus y$
- X-NOR (Equivalence) $f(x, y) = xy + \overline{x} \overline{y}$



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Some Terminology

 System specification: A description of the function of a system and of other characteristics required for its use.

- A function (table, algebraic) on a finite set of inputs.

- Double-rail logic*: both variables and their complements are considered as primary inputs.
- Single-rail logic: only variables are considered as primary inputs. Need inverters for their complements.

Gates and Combinational Networks

- Synthesis Procedure
- Example: Truth table for parity function on three variables

X	Y	Z	f
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

Synthesis Procedure

X	Y	Z	f
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

Minterm Canonical Form: $f(x, y, z) = \overline{x} \, \overline{y}z + \overline{x}y \, \overline{z} + x\overline{y} \, \overline{z} + xyz$

Two-level Gate Network

Minterm Canonical Form: $f(x, y, z) = \overline{x} \, \overline{y}z + \overline{x}y \, \overline{z} + x\overline{y} \, \overline{z} + xyz$



Incomplete Boolean Functions and Don't Care Conditions

Incomplete Boolean Functions and Don't Care Conditions

- n-variable incomplete Boolean function is represented by a truth table with n+1 columns and 2ⁿ rows.
- For those combinations of values in which a functional value is not to be specified, a symbol, --, is entered.
- The complement of an incomplete Boolean function is also an incomplete Boolean function having the same unspecified rows of the truth table.

Describing Incomplete Boolean Functions

X	Y	Z	F
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	
1	0	0	0
1	0	1	
1	1	0	0
1	1	1	1

Minterm canonical formula: $f(x, y, z) = \sum m(0, 1, 7) + dc(3, 5)$

Describing Incomplete Boolean Functions

X	Y	Z	F
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	
1	0	0	0
1	0	1	
1	1	0	0
1	1	1	1

Maxterm canonical formula: $f(x, y, z) = \Pi M(2,4,6) + dc(3,5)$

Describing Incomplete Boolean Functions

- Manipulating Boolean equations derived from incomplete Boolean functions is a very difficult task.
- In the next chapter, there are procedures for obtaining minimal expressions that can handle the don't care conditions.
- Can leverage don't care conditions to get simplified expressions for functions (smaller gate networks).

More on Don't Care Conditions

- <u>Satisfiability don't cares</u> of a subcircuit/subsystem consist of all input patterns that will never occur.
- <u>Observability don't cares</u> of a subcircuit/subsystem are the input patterns that represent situations when an output is not observed.



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Universal Gates

- A gate or set of gates is called universal if it can implement all Boolean functions.
- Standard universal gates:
 - AND, OR, NOT
 - Proof?
- Theorem:
 - A set of gates $\{G_1, \dots, G_n\}$ is universal if it can implement AND, OR, NOT.

NAND is Universal

- Recall $NAND(x, y) = \overline{x y} = \overline{x} + \overline{y}$
- NOT(x) =

 $\overline{x} = \overline{x} + \overline{x} = NAND(x, x)$

•
$$AND(x,y) =$$

 $xy = \overline{(xy)} = NOT((\overline{xy}))$
 $= NOT(NAND(x,y))$
 $= NAND(NAND(x,y), NAND(x,y))$

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$$OR(x, y) =$$

 $x + y = \overline{\overline{(x + y)}} = \overline{(\overline{x} \ \overline{y})}$
 $= NAND(NOT(x), NOT(y))$
 $= NAND(NAND(x, x), NAND(y, y))$

NOR is Universal

- Recall $NOR = \overline{x + y} = \overline{x} \, \overline{y}$
- NOT(x) = $\overline{x} = \overline{x} \, \overline{x} = NOR(x, x)$
- AND(x) = $xy = \overline{\overline{(xy)}} = \overline{(\overline{x} + \overline{y})}$ = NOR(NOT(x), NOT(y))= NOR(NOR(x, x), NOR(y, y))
- OR(x) = $x + y = \overline{\overline{(x + y)}} = NOT(\overline{x + y})$ = NOT(NOR(x, y))= NOR(NOR(x, y), NOR(x, y))

Class Exercise